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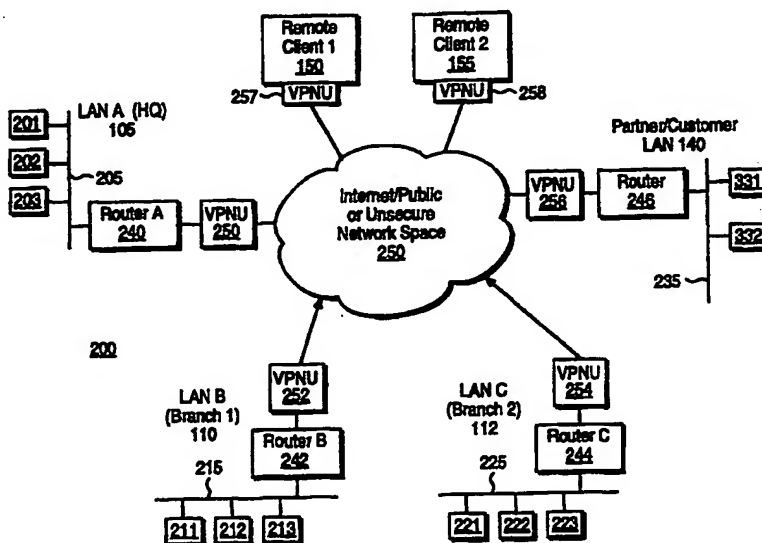
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| (71) Applicant: VPNET TECHNOLOGIES, INC. [US/US]; 1530 Meridian Avenue, San Jose, CA 95125 (US). | | | |
| (72) Inventor: GILBRECH, Sidney, A.; 1395 Richardson Avenue, Los Altos, CA 94024 (US). | | | |
| (74) Agents: THOMSON, William, E., Jr. et al.; McCutchen, Doyle, Brown & Enersen, Three Embarcadero Center, San Francisco, CA 94111 (US). | | | Published With international search report. |

(54) Title: AN APPARATUS FOR IMPLEMENTING VIRTUAL PRIVATE NETWORKS



(57) Abstract

Protocols and architecture for secure virtual private networks. Intraenterprise data communications are supported in a secure manner over the Internet or other public network space with the implementation of secure virtual private networks. Members of a virtual private network group exchange data that may be compressed, encrypted and authenticated, if the exchange is between members of the group.

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AN APPARATUS FOR IMPLEMENTING VIRTUAL PRIVATE NETWORKS

BACKGROUND OF THE INVENTION

1. Related Information

The present invention is related to the one described in Copending U.S. Patent application entitled "Architecture for Virtual Private Networks" Serial Number 08/874,090 [Attorney Docket No. 20155-701], assigned to the assignee of the present invention and filed concurrently herewith.

2. Field of the Invention

The present invention relates to the field of data communications. More particularly, the present invention relates to techniques for implementing secure virtual private networks over public or otherwise insecure data communications infrastructures.

3. Background

In recent years organizations have come to rely heavily on the ability to transmit electronic data between members of the organization. Such data typically includes electronic mail and file sharing or file transfer. In a centralized, single site organization, these transfers of electronic data are most commonly facilitated by a local area network (LAN) installed and operated by the particular enterprise.

Preventing unauthorized access to data traversing an enterprise's LAN is relatively straightforward. This applies to both unauthorized accesses by members of the enterprise and, more importantly, to third parties on the outside. As long as intelligent network management is maintained, unauthorized accesses to data traversing an enterprise's internal LAN are relatively easily avoided. It is when the enterprise spans multiple sites that security threats from the outside become a major concern.

For distributed enterprises that desire the conveniences of the above-described electronic data transfers, there are several options that exist today, but each with associated disadvantages. The first option is to interconnect the offices or various sites with dedicated, or private communications connections often referred to as leased lines. This is the traditional method organizations use to implement a wide area network (WAN). The disadvantages of implementing an enterprise owned and

controlled WAN are obvious: they are expensive, cumbersome and frequently underutilized if they are established to handle the peak capacity requirements of the enterprise. The obvious advantage to this approach is that the lines are dedicated for use by the enterprise and are therefore secure, or reasonably secure from eavesdropping or tampering by intermediate third parties.

An alternative to the use of dedicated communications lines in a wide area network is for an enterprise to handle intersite data distributions over the emerging public network space. Over recent years, the Internet has transitioned from being primarily a tool for scientists and academics to a mechanism for global communications with broad ranging business implications. The Internet provides electronic communications paths between millions of computers by interconnecting the various networks upon which those computers reside. It has become commonplace, even routine, for enterprises, even those in nontechnical fields, to provide Internet access to at least some portion of the computers within the enterprise. For many businesses this facilitates communications with customers, potential business partners as well as the distributed members of the organization.

Distributed enterprises have found that the Internet is a convenient tool to provide electronic communications between members of the enterprise. For example, two remote sites within the enterprise may each connect to the Internet through a local Internet Service Provider (ISP). This enables the various members of the enterprise to communicate with other sites on the Internet including those within their own organization. The limiting disadvantage of using the Internet for intra-enterprise communications is that the Internet is a public network space. The route by which data communication travel from point to point can vary on a per packet basis, and is essentially indeterminate. Further, the data protocols for transmitting information over the various networks of the Internet are widely known, and leave electronic communications susceptible to interception and eavesdropping with packets being replicated at most intermediate hops. An even greater concern arises when it is realized that communications can be modified in transit or even initiated by impostors. With these disconcerting risks, most enterprises are unwilling to subject their proprietary and confidential internal communications to the exposure of the

public network space. For many organizations it is common today to not only have Internet access provided at each site, but also to maintain the existing dedicated communications paths for internal enterprise communications, with all of the attendant disadvantages described above.

5 While various encryption and other protection mechanisms have been developed for data communications, none completely and adequately addresses the concerns raised for allowing an enterprise to truly rely on the public network space for secure intra-enterprise data communications. It would be desirable, and is therefore an object of the present invention to provide such mechanisms which would allow the distributed enterprise to rely solely on the public network space for intra-enterprise communications without concern for security risks that presently exist.

SUMMARY OF THE INVENTION

From the foregoing it can be seen that it would be desirable and advantageous to develop protocols and architecture to allow a single organization or enterprise to rely on the public network space for secure intraorganizational electronic data communications. The present invention is thus directed toward the protocols and architecture for implementing secure virtual private networks over the Internet or other public network apparatus. The architecture of the present invention introduces a site protector or virtual private network (VPN) unit which moderates data communications between members of a defined VPN group. In accordance with one embodiment of the present invention, the site protector resides on the WAN side of the site's router or routing apparatus which is used to connect the enterprise site to the Internet. In alternative embodiments, the site protector will reside on the LAN side of the router. The essential point for all embodiments is that the site protector be in the path or all relevant data traffic.

25 To ensure secure data communications between members of the same VPN group, the site protector or VPN Unit implements a combination of techniques for data packet handling when packets are to be sent between members of the group. The packet handling processes include various combinations of compression, encryption and authentication, the rules for each of which may vary for members of different groups. For each group defined as a virtual private network, the various parameters

defining the compression, encryption and authentication are maintained in lookup tables in the associated VPN Units. The lookup tables maintain information not only for fixed address members of the group but support is also provided for remote clients. This ability allows remote users to dial into a local Internet Service Provider and still maintain membership in a virtual private network group for secure communications over the Internet with other members of the group. In the case of a remote client, the site protector may, in one embodiment, be simulated by software running on the remote client.

In other aspects of the present invention, the VPN Units or site protectors may be dynamically configured to add or subtract members from the virtual private network group or recognize their movement, or change other parameters affecting the group. Various other packet handling aspects of the invention include addressing the problem of some data packets growing too large by the inclusion of encryption and authentication information. Another packet handling aspect provides a mechanism for Internet communications which hides information identifying the source and destination of the data packet. In this aspect of the present invention, the VPN Units are treated as the source and destination for the Internet communication data packets with the VPN Units encapsulating the source and destination addresses of the endstations.

There is also disclosed a hardware architecture and implementation for a VPN Unit. This embodiment is one designed to reside on the WAN side of a given site's router. In the illustrated embodiment, a combination of computer hardware and software are provided for carrying out the processes of the VPN apparatus for compression, encryption and authentication under the direction of a microprocessor.

BRIEF DESCRIPTION OF THE DRAWINGS

The objects, features and advantages of the present invention will be apparent from the following detailed description, in which:

Figure 1 illustrates a prior art configuration for an exemplary enterprise's intraenterprise communication architecture.

Figure 2 illustrates an enterprise communication scenario in accordance with the present invention utilizing the Internet or other public network space as the vehicle for conveying messages between members of a virtual private network.

Figure 3 illustrates a flow diagram for the handling of a packet being
5 transmitted from one member of a virtual private network group to another member over the Internet.

Figure 4 illustrates the handling of a data packet received over the Internet by one member of a virtual private network group from another member.

Figure 5 illustrates graphically the life cycle of a data packet being sent from
10 one member of a virtual private network group to another over the Internet.

Figure 6 illustrates an alternate life cycle of a data packet being sent from one member of a virtual private network group to another over the Internet where the source and destination addresses of the group members are also concealed.

Figure 7 illustrates an architectural block diagram for an implementation of a
15 virtual private network unit in accordance with one embodiment of the present invention.

Figure 8 illustrates a more detailed block diagram for one implementation of a virtual private network unit in accordance with one embodiment of the present
invention.

20 DETAILED DESCRIPTION OF THE INVENTION

An apparatus and apparatus are disclosed for implementing secure virtual private networks for enterprise communications over the Internet or other public network space. Although the present invention is described predominantly in terms of utilizing the Internet as a communications medium, the concepts and methods are
25 broad enough to accomplish the implementation of secure virtual private networks over other public or insecure communications media. Throughout this detailed description, numerous specific details are set forth such as particular encryption or key management protocols, in order to provide a thorough understanding of the present invention. To one skilled in the art, however, it will be understood that the present
30 invention may be practiced without such specific details. In other instances, well-

known control structures and apparatus components have not been shown in detail in order not to obscure the present invention.

In many instances, components implemented by the present invention are described at an architectural, functional level. Many of the elements may be configured using well-known structures, particularly those designated as relating to various compression or encryption techniques. Additionally, for logic to be included within the apparatus of the present invention, functionality and flow diagrams are described in such a manner that those of ordinary skill in the art will be able to implement the particular methods without undue experimentation. It should also be understood that the techniques of the present invention may be implemented using a variety of technologies. For example, the virtual private network unit or site protector to be described further herein may be implemented in software running on a computer apparatus, or implemented in hardware utilizing either a combination of microprocessors or other specially designed application specific integrated circuits, programmable logic devices, or various combinations thereof. It will be understood by those skilled in the art that the present invention is not limited to any one particular implementation technique and those of ordinary skill in the art, once the functionality to be carried out by such components is described, will be able to implement the invention with various technologies without undue experimentation.

Referring now to Figure 1 there is shown a traditional scenario for intra-enterprise data communications for a distributed organization. In this illustration of an exemplary organization configuration, the enterprise consists of a headquarters location 105 with additional sites or branches 110 and 112, respectively. In modern organizations, such as the exemplary one of Figure 1, the headquarters' site 105 as well as the branch sites 110 and 112 may each comprise numerous personnel, many of whom are provided with computers or work stations with network access. The internal network configurations at the headquarters for branches may take many forms including one or several local area networks (LANs). For intersite communications between headquarters and the branches, dedicated or leased communications lines 115 and 120 may be provided. In addition, an optional dedicated communications path 125 may be provided between the branches 110 and 112. As an alternative to the

optional dedicated communications line 125 between the branches, data packets between branch 110 and branch 112 may be routed through the headquarters' network equipment.

In addition to the dedicated communications lines between the headquarters and the various branches, it is common today to provide computer users within an organization access to the Internet for electronic mail to external parties as well as for doing various types of research over the Internet using such tools as the World Wide Web, etc. As shown in Figure 1, the usual scenario where the headquarters' site 105 and the branches 110 and 112 are each separately provided with direct access to Internet Service Providers 130, 133 and 136, respectively. This facilities the users at the various sites with their access to the Internet for the above purposes. In an alternate configuration, it may be that only the headquarters site 105 is provided with access to an Internet service provider 130 and that users of the computers of the branch sites 110 and 112 will connect to the Internet through headquarters via their dedicated communications paths 115 and 120. The downside to this alternate configuration is that it greatly increases the bandwidth utilization on the dedicated lines, perhaps to the point of saturation. An advantage is that only one gateway to the Internet need be provided for the organization which simplifies enforcing security constraints on connections to the outside world.

In the exemplary organization 100, it is also shown that in some circumstances it may be desirable to allow customers or other business partners to dial in directly to the computer network of the organization. In Figure 1 it is illustrated that the customer 140 may in fact carry out such communications over a communications path 145 which may be a dedicated line provided between the customer and the organization for the customer's convenience. The path 145 may also be a dial-up line which the customer might use only sporadically. Consistent with the emerging use of the Internet and its popularity, the customer 140 is shown having its own Internet connection through ISP 148.

Finally, there is shown in Figure 1 that it is frequently desirable for other members of the enterprise who may be on the road or working from home or other remote locations to exchange data with other members of the enterprise. There is thus

shown remote clients 150 and 155 communicating with the headquarters over long distance telephone lines 157 and 158. This example assumes that the remote clients are in a truly remote location from the headquarters. The remote clients 150 and 155 are also respectively shown having local access to the Internet through local ISPs 160 and 165.

The above description of an enterprises data communications configuration according to Figure 1 illustrates the disadvantages described in the previous section. These disadvantages are eliminated by implementation of the present invention as illustrated generally with reference to Figure 2. In the enterprise network communication configuration 200 illustrated in Figure 2, the headquarters 105, first branch 110 and second branch 112 of the organization are illustrated in a more detailed logical way then presented in Figure 1. Thus, the headquarters 105 is illustrated with three endstations 201, 202 and 203, respectively coupled to communicate data packets over local area network (LAN) 205. Likewise, the branch site 110 is shown having a plurality of endstations 211, 212 and 213 respectively coupled to communicate data locally over LAN 215. Finally, the second branch site 112 is shown with an illustrative set of computer stations 221, 222 and 223 connected to communicate over LAN 225. The customer site 140 is also illustrated in Figure 2 as comprising of plurality of computers illustrated by 331 and 332 coupled to communicate over the customer's LAN 235. The local area networks utilized for data communications within the headquarters, customer and branch sites may adhere to a wide variety of network protocols, the most common of which are Ethernet and Token Ring.

As can be seen in Figure 2, the dedicated communications lines between the headquarters site 105 and the branch sites 110 and 112 as well as between the headquarters site 105 and the customers site 140 have been eliminated. Instead, in accordance with the present invention data communications between members of the organization are intended to be carried out over the Internet or other public network space. For purposes of the present invention, it will be assumed that it is the widely emerging Internet that will be the medium for data packet transfers between members of the organization.

Each of the LANs for the particular sites illustrated in Figure 2 ultimately interconnect to the Internet 250 through an associated routing or gateway device which are identified as routers 240, 242, 244 and 246, respectively. It is to be understood that data packets conveyed between a various sites illustrated in 200
5 would traverse, in many cases, a plurality of additional routing devices on their way between the source and destination sites for the packets. The mechanisms for data packet transfers over the Internet are well known and are not described in great detail herein. It is understood that data packets are assembled in accordance with the Internet Protocol (IP) and are referred to herein as IP packets regardless of the version
10 of the Internet protocol presently in effect. In the case of the remote clients 150 and 155 illustrated in Figure 2 it is understood that they utilize communication software to dial up a local Internet service provider which itself provides the gateways necessary for communications over the Internet 250.

As has been described above, prior efforts to utilize the Internet for secure data
15 communications have required an awareness or implementation of security considerations at the endstations. This is disadvantageous when transparency to an end user is desirable. The present invention, on the other hand is transparent to end users with data communications over the Internet occurring exactly as they appear to have before. However, for users identified as members of the same virtual private
20 network, data communications are handled in a manner that assures the security and integrity of the data packets. Illustrated in Figure 2, between the Internet 250 and each of the respective routers 240, 242, 244 and 246, are Virtual Private Network Units (VPN Units) 250, 252, 254 and 256. In accordance with the particular illustrated embodiment of the present invention, the VPN Units reside between a site's
25 router and the path to the Internet. It should be understood that this placement of VPNU units in the overall system architecture represents only one placement choice. It will be clear from the materials that follow that the key point with respect to VPNU placement is that they reside in the path of data traffic. In many embodiments, it may in fact prove desirable to situate the VPNU on the LAN side of a site's router. As will
30 be described in more detail below, the VPN Units maintain lookup tables for identifying members of specific virtual private network groups.

When a data packet is sent between source and destination addresses that are both members of the same VPN group, the VPN Unit will process the data packet from the sending side in such a way as to ensure that it encrypted, authenticated and optionally compressed. Likewise, the VPN Unit servicing the site where the
5 destination address is located will detect that a packet is being propagated between members of the same VPN group. The receiving VPN Unit will handle the process of decrypting and authenticating the packet before forwarding it toward the destination endstation. In this way, secure data communications between end users is effected in a manner that is transparent to the end users. In the case of remote clients 150 and
10 155, the VPN Unit may be simulated in software which operates in conjunction with the communication software for connecting the remote client to the associated local Internet service provider.

The functionality of the VPN Units will be described with reference to the following figures beginning with the flowchart of Figure 3. When a data packet
15 originates from an endstation, such as endstation 202 of LAN 205 at site 105, and its destination is to a remote site, other than the headquarters site 105, it will initially be treated as an ordinary Internet data packet transfer. The packet will proceed from the endstation 202 over the LAN 205 to the routing device 240 which will encapsulate the data packet in accordance with the Internet Protocol, forming an outbound IP packet.
20 On its way out of the site, the IP packet will pass through the associated VPN Unit for the site. The flowchart illustrated at Figure 3 shows the functional operation of a VPN Unit for an outbound packet that is received thereby. The Transmit Packet procedure 300 begins when the outbound data packet is received at the VPN Unit at step 310. At decision box 320, it is determined whether or not the source and destination addresses
25 for the data packet are both members of the same VPN group. This determination may be made with reference to lookup tables that are maintained by the VPN Units or reference to other memory mechanisms. This step may be thought of as member filtering for data packets being transmitted between the particular site and the VPN Unit which services it. If the source and destination address for the data packet are
30 not both members of the same VPN group, then at step 330 the packet is forwarded to the Internet as ordinary Internet traffic from the site as though the VPN Unit were not

involved. In which case, the procedure ends at step 335. In one alternative embodiment, it may be desirable to discard data traffic that is not destined between members of a VPN group rather than forwarding it as unsecure traffic. In another alternative embodiment, it may be desirable to provide the option to either pass or
5 discard non-VPN-group data traffic.

If, at decision box 320, the member filter, it is determined that both the source and destination addresses for the data packet are members of the same VPN group, then the data packet is processed at step 340 undergoing various combinations of compression, encryption and authentication. The lookup tables maintained by the
10 VPN Unit 250 and all of the VPN Units, in addition to identifying members of particular VPN groups, also identify whether or not data packets transferred between members of the particular VPN group are to be compressed and if so, what algorithm is to be used for compression. Many possible compression algorithms are well-known, but in one embodiment of the invention, LZW compression is implemented.
15 The lookup table for the VPN group of which the source and destination addresses are members also identifies the particular encryption algorithm to be used for data packets traversing the Internet for that VPN group as well as the authentication and key management protocol information to be used thereby. As an alternative to lookup tables, the VPN Unit may be programmed to always use the same algorithms for all
20 VPN groups.

The particular packet processing algorithms to be used for VPN traffic may vary, so long as the lookup tables in both the sending and receiving VPN Units identify the same compression, encryption and authentication rules and are capable of implementing and deimplementing them for members of the same group. It is to be
25 understood that a single VPN Unit may serve multiple VPN groups and that particular addresses may be members of multiple groups. Thus, at step 340, when a packet is destined from one member of the VPN group to another, the packet is processed according to the compression, encryption and authentication rules identified in the VPN Unit tables for that particular VPN group. Then, at step 350, the processed
30 packet is forwarded toward the destination address over the Internet. The procedure of the sending VPN Unit then ends at step 355.

The receiving VPN Unit reverses the above processes for VPN traffic as illustrated by the flowchart of Figure 4. The Receive Packet procedure 400 begins at step 410 when an inbound data packet is received from the Internet at the receiving VPN Unit. At decision box 420, the inbound data packet is examined to determine if the source and destination addresses of the data packet are both members of the same VPN group. It is assumed that the lookup tables maintained by all of the VPN Units are both consistent and coherent. If the inbound data packet is determined not to be VPN traffic, then the packet is passed through and forwarded to the receiving site as though it were normal Internet data traffic at step 430. In which case the process ends at step 435. In one alternative embodiment it may be desirable to discard incoming data traffic that is not from an identified member of a VPN group supported by the VPNU.

For data packets that are determined to be VPN traffic at decision box 420, the VPN Unit will process the inbound packet to recover the original data packet as it was provided from the source endstation. The lookup table maintained by the receiving VPN Unit will identify the compression, encryption and authentication rules used for the VPN group and reconstruct the original IP packet in accordance with those rules at step 440. Then, the reconstructed packet will be delivered to the site of the destination address at 450 with the procedure ending at step 455.

Figure 5 illustrates graphically the life cycle of the data packet sent between two members of the same VPN group. The data packet originates from a source 500 and propagates from the sources site through its associated router to generate IP data packet 510. The data packet 510 is not intended to illustrate all the fields associated with a complete IP data packet, but shows the relevant portions for this discussion which include the destination address, source address and the payload information of the packet. The data packet 510 is then examined by the VPN Unit which determines whether the data packet is traffic between members of an identified VPN group. The VPN Unit 520 processes the packet in accordance with the packet processing procedures described above with respect to Figure 3 with the resulting packet being illustrated as packet 530. Packet 530 still identifies the destination and source

addresses of the data packet, but the remainder of the packet is encrypted, and optionally compressed.

Following processing by the outbound VPN Unit, the data packet is propagated through the Internet to 550 with the destination and source information identifying to the associated routers of the Internet the path by which the packet should ultimately take to reach its destination. The packet emerges from the Internet at the edge of the destination site as data packet 540 which is essentially identical to the data packet 530. The packet is "deprocessed" by the receiving VPN Unit 550 which restores the original packet into its form 560 for delivery to the ultimate destination through the receiving site's associated router at destination 570.

As was described above, the present invention approach to virtual private networks supports not only optional compression of data packets, but encryption and authentication techniques as well. One emerging standard for key management in connection with Internet Protocol data transfers with authentication is referred to as simple key management for Internet Protocol (SKIP) which is described by US Patent 5,588,060 assigned to Sun Microsystems, Inc. of Mountain View, CA. Authenticated data transfers using SKIP support a mode of data transfer referred to as tunnel mode. The above described data transfer with respect to Figure 5 illustrates a transport mode of operation in which the data and source addresses are exposed as the data packet traverses the Internet. In tunnel mode, an added measure of security may be provided by encapsulating the entire data packet in another packet which identifies the source and destination addresses only for the VPN Units. This conceals the ultimate source and destination addresses in transit.

Figure 6 illustrates the life cycle of a data packet being propagated from a source 600 to a destination 670 utilizing tunnel mode. In this mode of operation, the data packet 610 is processed by outbound VPN Unit 620 which generates a resulting packet 630. The resulting packet 630 encrypts and compresses (optionally) not only the data payload of the packet, but the destination and source addresses of the endstations as well. The encapsulated packet is then provided with an additional header that identifies that the source of the packet is the outbound VPN Unit 620 and that the destination is the inbound VPN Unit 650. Thus, the packet 640 which

emerges from the Internet is identical to the packet 630 with respect to its source and address information and encapsulated payload. The packet is decomposed by the inbound VPN Unit 650 to reconstruct the original data packet at 660 for delivery to the destination 670.

5 The overall architecture of the present invention is robust. It allows end users the convenience of proprietary data communications to take place over a public network space such as the Internet. The architecture of the present invention also allows a wide variety of compression, encryption and authentication technologies to be implemented, so long as the VPN Units at each end of the transaction support the
10 associated protocols. The present invention is also capable of working in concert with traditional Internet security mechanisms such as corporate firewalls. A firewall might operate in series with the VPN Unit at a given site, or, intelligently be configured in a single box with the VPN Unit to provide parallel firewall and VPN Unit security functions.

15 **Architecture for a Virtual Private Network Unit**

 The above discussion relates to the functionality for implementing virtual private networks. There is now described a hardware architecture and implementation for a virtual private network unit in accordance with one embodiment of the present invention. Referring now to Figure 7 there is illustrated a general block diagram for a
20 virtual private network unit 700. The general architecture of the VPN Unit 700 is essentially that of a specially tailored personal computer (PC) apparatus architecture. The overall operation of the VPN Unit 700 is driven by a central processing unit (CPU) 705 which is coupled to communicate with the other components of the apparatus over a apparatus bus 710. In the general architecture illustrated for VPN
25 Unit 700, the memory apparatus 720 is also illustrated as residing on the apparatus bus 710. Those of ordinary skill in the art will of course recognize that various memory configurations may be implemented, some of which may include a dedicated memory bus for higher speed throughput to the memory apparatus 720.

 The VPN Unit 700 is designed to be an intermediate device between a site's
30 local area network and the Internet or other public network space. Thus, the VPN Unit 700 is illustrated as having an I/O controller 730 for coupling the VPN Unit 700

to the Internet or other public network space. Similarly, an I/O controller 740 is provided for coupling the VPN Unit 700 to the site's edge router. In accordance with the above-described functional requirements, the VPN Unit 700 resides between the site's router and the Internet. In this implementation, the I/O controller 740 is responsible for data packet transmissions between the router for the site and the VPN Unit 700. In an alternate embodiment of the present invention, the VPN Unit 700 may be disposed between a site's local area network and its router, in which case the I/O controller 730 would be responsible for data packet transmissions between the VPN Unit 700 and the router, while the I/O controller 740 would effectively be the VPN Unit's interface to the local area network. In this alternative embodiment, data packet transmissions would probably have to be handled in accordance with the network protocol standard of the local area network rather than the public network space IP transmissions as described above. In either event, the general architecture for VPN Unit 700 illustrates that the two I/O controllers 730 and 740 are coupled to communicate over the apparatus bus 710 through I/O control logic 750. In accordance with this architecture, the I/O control logic 750 is responsible for arbitrating access to the apparatus bus for the data packets coming in and going out, and for assuring the necessary throughput for the data connections.

As has been described, when a data packet is received at the VPN Unit 700, whether outbound or inbound, a determination is made by examining the data packet's header as to whether or not both the source and destination addresses represent members of the same VPN group. The determination with respect to one embodiment requires interrogation of a look-up table, which maintains the identities of the various VPN groups serviced by the VPN Unit 700, as well as the group members and the various processing parameters for data packets that are exchanged between group members. In accordance with the general architecture of the VPN Unit 700, this interrogation may be handled by the CPU 705 with reference to look-up tables that would be maintained in the memory apparatus 720.

For data packets that are to be processed as packets being sent between members of a supported VPN group, it has been described that compression and decompression functions as well as encryption and decryption functions need to be

performed for such data. Thus, the VPN Unit 700 architecture includes a compression engine 760, coupled to communicate with the rest of the apparatus over the apparatus bus 710, and an encryption engine 770 similarly coupled to the apparatus bus 710.

For supporting functions such as the various key management protocols described
5 above which result in the generation of packet encapsulation headers, processing may be carried out by the CPU 705 with packets being assembled within the CPU 705 or at specified memory locations in the memory apparatus 720. Alternatively, logic may be incorporated in the VPN Unit 700 designed specifically for the support of such key management protocols or other packet processing operations.

10 There is finally shown in the VPN Unit 700 architecture that optional I/O devices 780 may be interfaced to the VPN Unit 700 through I/O control logic 790. This may facilitate management of the unit either by allowing direct keyboard access to control the processor, or provide a connection path for other communications devices which may need to communicate with the VPN Unit 700.

15 Referring now to Figure 8, there is illustrated a more detailed block diagram for one implementation of a VPN Unit in accordance with one embodiment of the present invention. In the implemented embodiment of VPN Unit 800 an architecture is disclosed which is driven by microprocessor 805, a 100 MHz Intel 486 DX4. The apparatus bus for this configuration is the VSIA bus 810 to which the other
20 components of the apparatus are coupled. The memory apparatus in this case is supplied by several banks of fast dynamic random access memory (DRAM) 820. In accordance with this implemented embodiment, a PC chip set implementation from the Opti Corporation is provided at 825 for controlling the overall operation of the apparatus. The Opti chip set, as well as providing a apparatus controller functionality
25 for the 486 DX4 microprocessor, may also be utilized to provide an interface to an optional ISA bus 828 to which other peripherals may be coupled to the apparatus.

In the implemented VPN Unit 800, the I/O connections to the public network space 830 are provided by a series of multi-protocol transceivers with a DB-25 connector. Similarly, the connections for the private network side of the unit are
30 provided through additional multi-protocol transceivers and an additional DB-25 connector. The I/O connections to the public and private networks are handled

through a dual-port SCC 845 which includes a plurality of input and output buffers for transmission and receipt of data packets to the VPN Unit 800. The data packet I/O controller is coupled to the VSIA bus 810 through FPGA Control and Glue Logic 850.

5 The compression engine for the implemented VPN Unit 800 is provided by an implementation of the STAC compression algorithm using the STAC chip 9710 compression engine 860, which is coupled to a dedicated bank of SRAM 865 to support its processing. Finally, in the implemented embodiment the encryption engine is provided by a DES engine 870, which may be implemented in either an off-the-shelf special integrated circuit configuration or one designed to operate in
10 accordance with the operation of the VPN Unit 800. In accordance with an alternative embodiment of the present invention, it is conceivable that the compression engine and encryption engine could be handled by software routines running on the microprocessor 805 or other general-purpose processing logic. Other alternative embodiments may use later generations of microprocessors which operate at higher
15 speeds and may require implementation of a different apparatus bus such as the emerging PCI bus standard. Those of ordinary skill in the art will recognize alternative and various embodiments which may be designed to suit a particular purpose.

There has thus been described a protocol and architecture for implementing
20 virtual private networks for using a public network space for secure private network data communications. Although the present invention has been described with respect to certain exemplary and implemented embodiments, it should be understood that those of ordinary skill in the art will readily appreciate various alternatives to the present invention. Accordingly, the spirit and scope of the present invention should
25 be measured by the terms of the claims which follow.

CLAIMS

What is claimed is:

- 1 1. An apparatus for providing secured data communications between
2 members of a virtual private network group comprising:
3 input/output (I/O) circuitry for receiving and transmitting data packets
4 between the members of said virtual private network group;
5 a system bus in communication with said I/O circuitry for conveying data
6 between components of said system
7 a compression engine in communication with said system bus for compressing
8 outbound data packets and decompressing inbound data packets;
9 an encryption engine in communication with said system bus for encrypting
10 outbound data packets and decrypting inbound data packets;
11 a central processing unit (CPU) in communication with said system bus for
12 controlling the processing of data packets by said system, said CPU for determining
13 which data packets are being sent between members of said virtual private network
14 and for determining the data flow for said data packets in accordance with
15 predetermined parameters for said virtual private network; and
16 a memory device in communication with said system bus for maintaining a list
17 of members of said virtual private network and for storing said predetermined
18 parameters.
- 1 2. The apparatus of claim 1 wherein said I/O circuitry comprises:
2 a private I/O port for coupling said apparatus to a site's private network;
3 a public I/O port for coupling said apparatus to a public network space; and
4 I/O control logic in communication with said system bus and coupled to said
5 private and public I/O ports for controlling data packet flow between said apparatus
6 and said members of said virtual private network.
- 1 3. The apparatus of claim 2 wherein said CPU generates encapsulation
2 headers for outbound data packets in accordance with a key management protocol.
- 1 4. The apparatus of claim 3 wherein said key management protocol
2 comprises the Simple Key Management for Internet Protocol (SKIP).

1 5. The apparatus of claim 2 wherein said encryption engine comprises an
2 application specific integrated circuit for performing DES encryption.

1 6. The apparatus of claim 2 wherein said encryption engine comprises an
2 application specific integrated circuit for performing DES encryption or triple-DES
3 encryption in accordance with said predetermined parameters.

1 7. The apparatus of claim 1 wherein said compression engine comprises
2 an integrated circuit for performing LZW compression.

1 8. The apparatus of claim 1 wherein said memory comprises lookup
2 tables for identifying all the virtual private network groups supported by said
3 apparatus and the members of said group wherein said members are each identified by
4 a network address and wherein a single network address may identify a member of
5 multiple groups.

1 9. An apparatus for securely exchanging data packets between members
2 of a virtual private network group comprising:

3 a first computer at a first site, said first computer having a first network
4 address;

5 a first router associated with said first site for routing data packets
6 originating from said first computer over a public network;

7 a first virtual private network unit disposed between said router and
8 said public network, said first virtual public network unit for identifying virtual
9 private network group data traffic and for securing said data traffic by manipulating
10 said data traffic according to packet manipulation rules maintained by said virtual
11 private network unit;

12 a second router associated with a second site for coupling said second
13 site to the public network;

14 a second virtual private network unit disposed between said second
15 router and the public network for intercepting network traffic destined for said second
16 site, said second virtual public network unit for detecting virtual private network
17 group traffic and for recovering original packet data; and

18 a second computer at said second site, said second computer having a second
19 network address for receiving said packet data, wherein said first and second
20 virtual private network units respectively comprise:
21 input/output (I/O) circuitry for receiving and transmitting data packets
22 between the members of said virtual private network group;
23 a system bus in communication with said I/O circuitry for conveying data
24 between components of said apparatus;
25 a compression engine in communication with said system bus for compressing
26 outbound data packets and decompressing inbound data packets;
27 an encryption engine in communication with said system bus for encrypting
28 outbound data packets and decrypting inbound data packets;
29 a central processing unit (CPU) in communication with said system bus for
30 controlling the processing of data packets by said apparatus, said CPU for determining
31 which data packets are being sent between members of said virtual private network
32 and for determining the data flow for said data packets in accordance with
33 predetermined parameters for said virtual private network; and
34 a memory device in communication with said system bus for maintaining a list
35 of members of said virtual private network and for storing said predetermined
36 parameters.

1 10. The system of claim 9 wherein said I/O circuitry comprises:
2 a private I/O port for coupling said apparatus to a site's private network;
3 a public I/O port for coupling said apparatus to a public network space; and
4 I/O control logic in communication with said system bus and coupled to said
5 private and public I/O ports for controlling data packet flow between said apparatus
6 and said members of said virtual private network.

1 11. The system of claim 10 wherein said CPU generates encapsulation
2 headers for outbound data packets in accordance with a key management protocol.

1 12. The system of claim 11 wherein said key management protocol
2 comprises the Simple Key Management for Internet Protocol (SKIP).

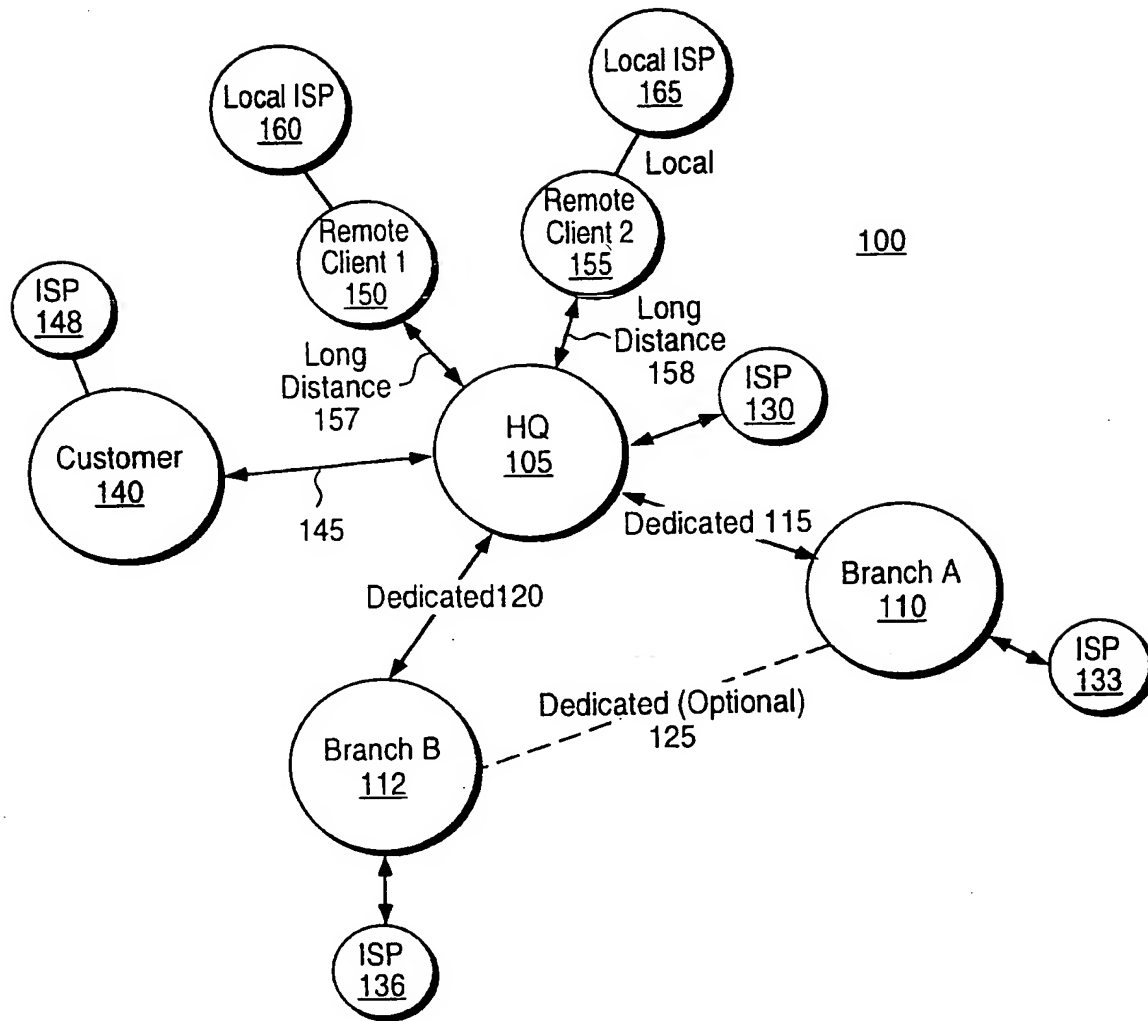
1 13. The system of claim 10 wherein said encryption engine comprises an
2 application specific integrated circuit for performing DES encryption.

1 14. The system of claim 10 wherein said encryption engine comprises an
2 application specific integrated circuit for performing DES encryption or triple-DES
3 encryption in accordance with said predetermined parameters.

1 15. The system of claim 10 wherein said compression engine comprises an
2 integrated circuit for performing LZW compression.

1 16. A system of claim 10 wherein said memory comprises lookup tables
2 for identifying all the virtual private network groups supported by said apparatus and
3 the members of said group wherein said members are each identified by a network
4 address and wherein a single network address may identify a member of multiple
5 groups.

FIG. 1 (Prior Art)



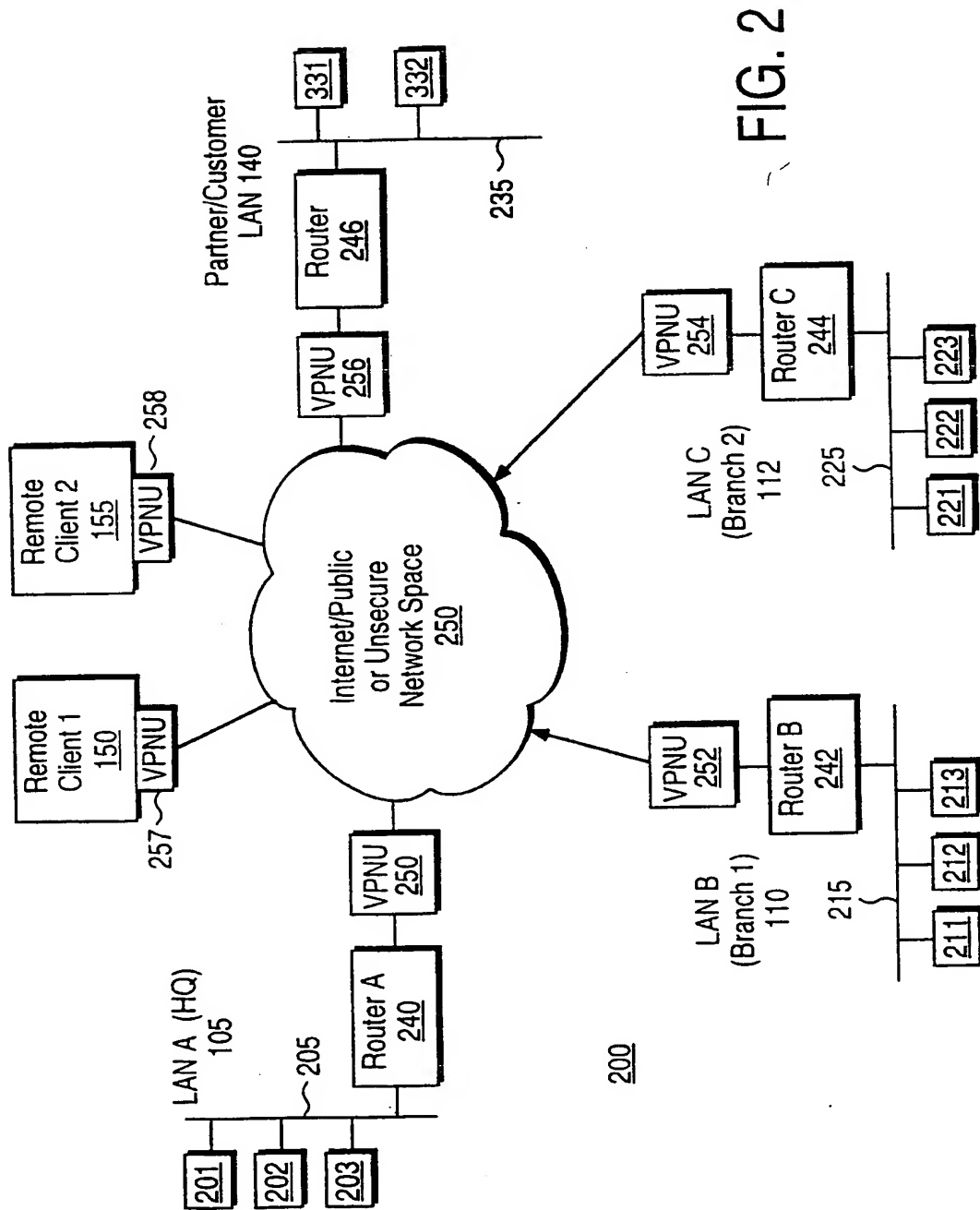


FIG. 2

FIG. 3

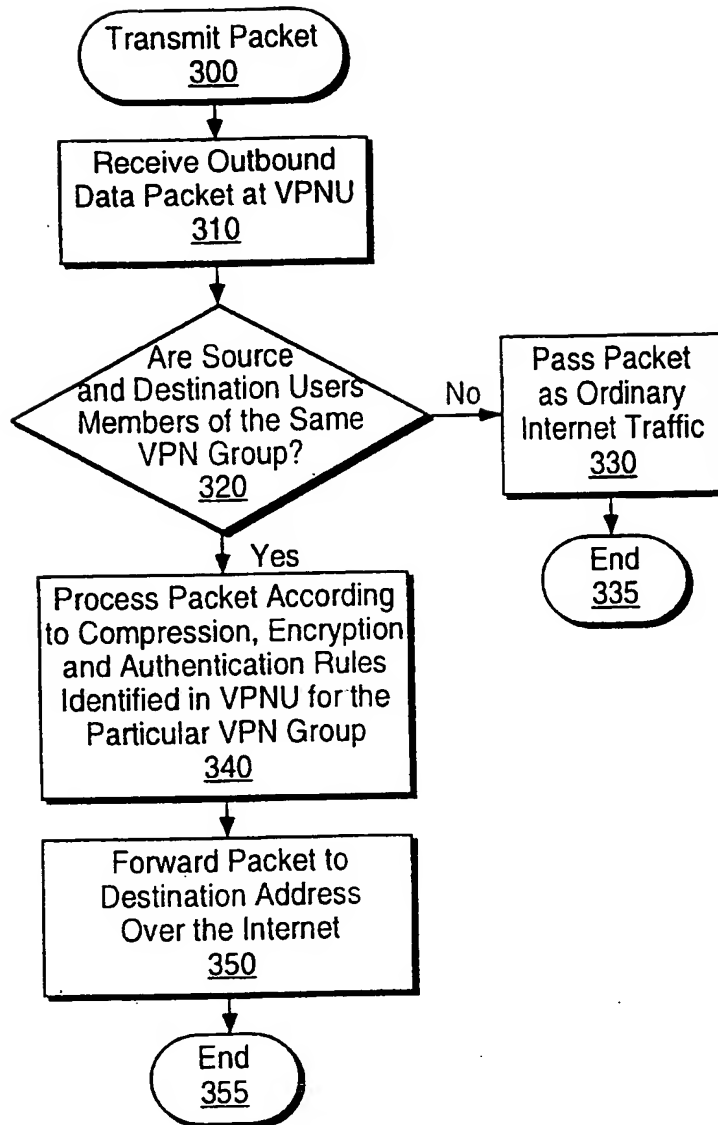


FIG. 4

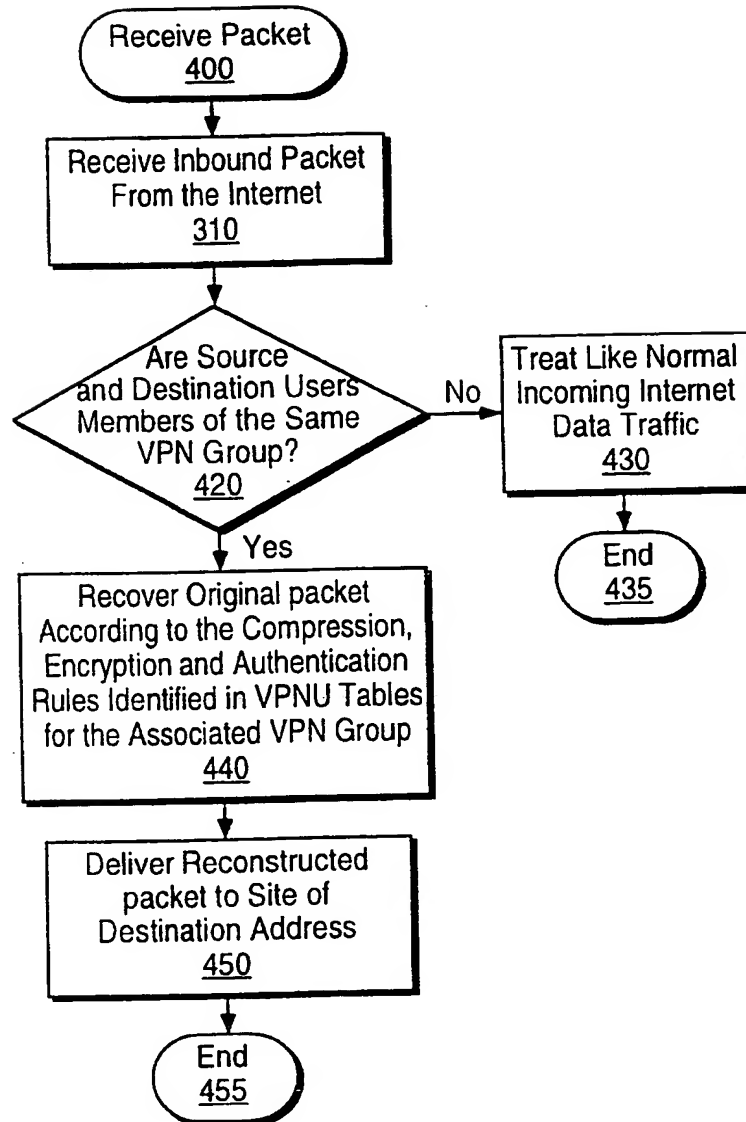


FIG. 5

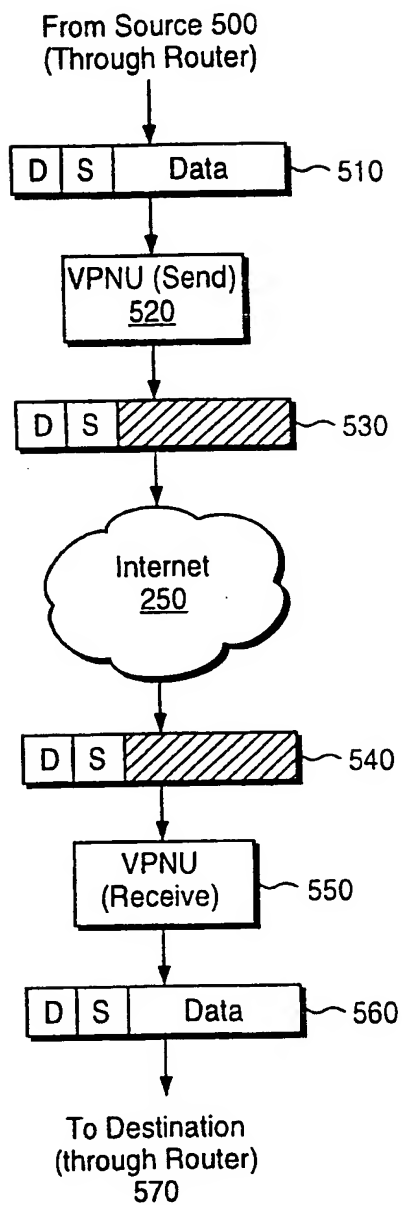
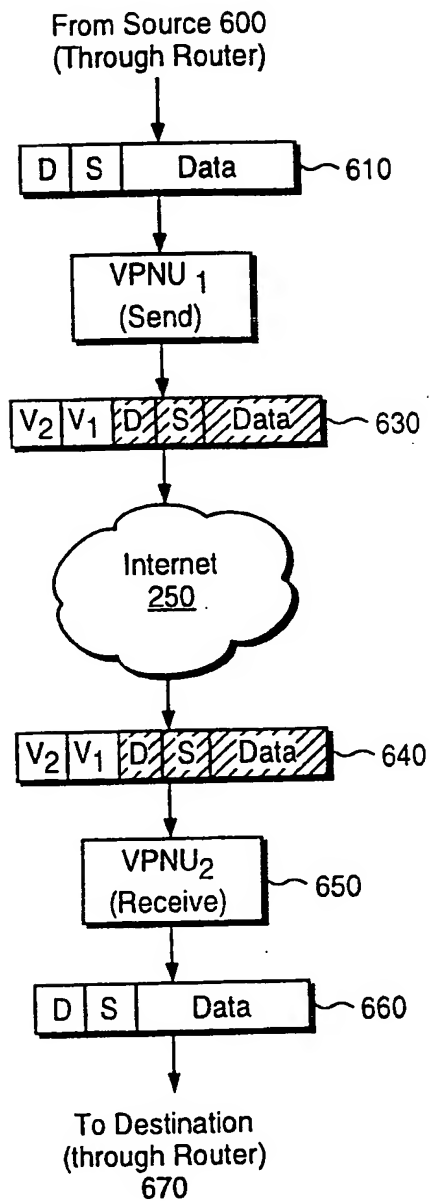


FIG. 6



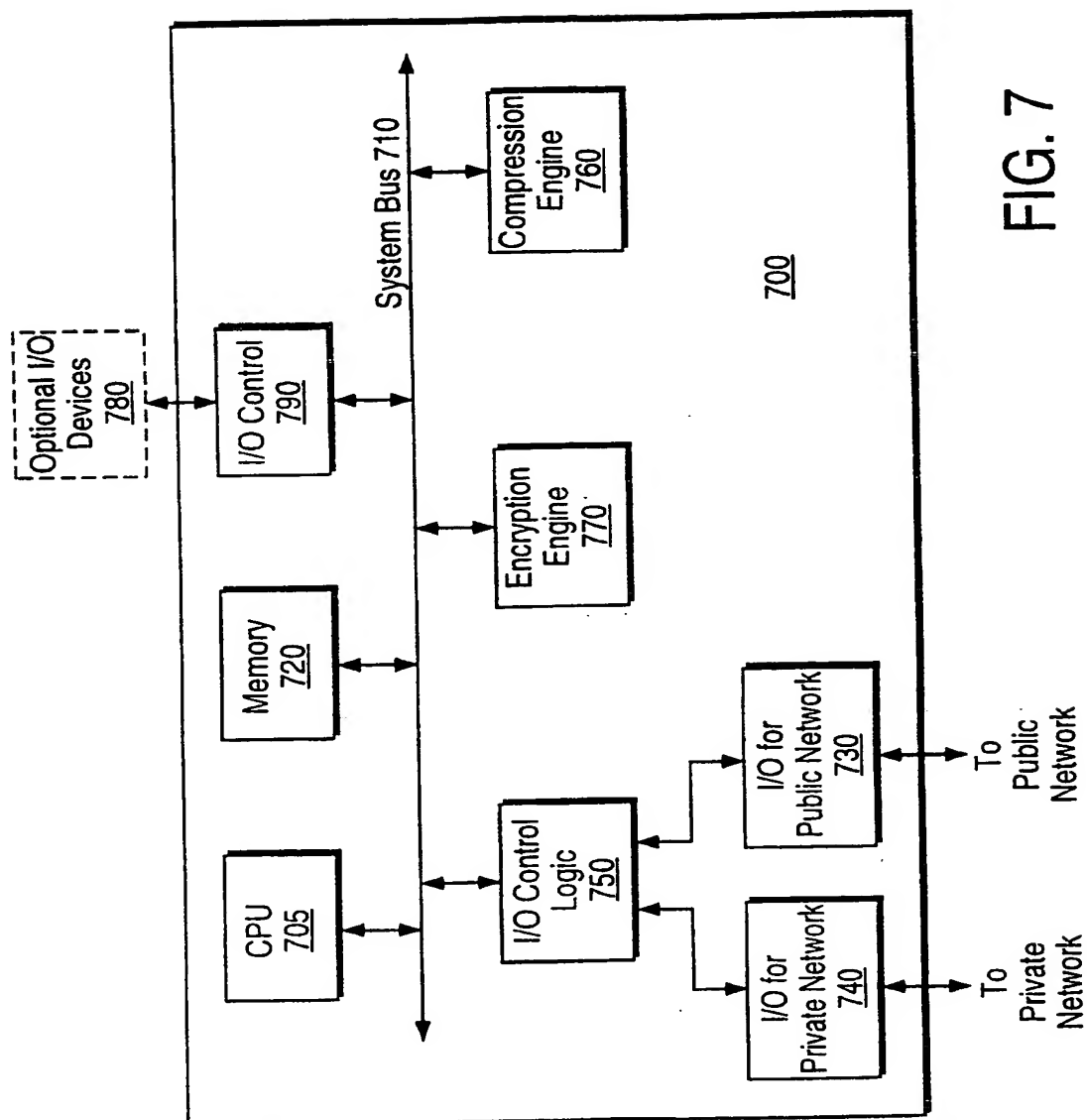


FIG. 7

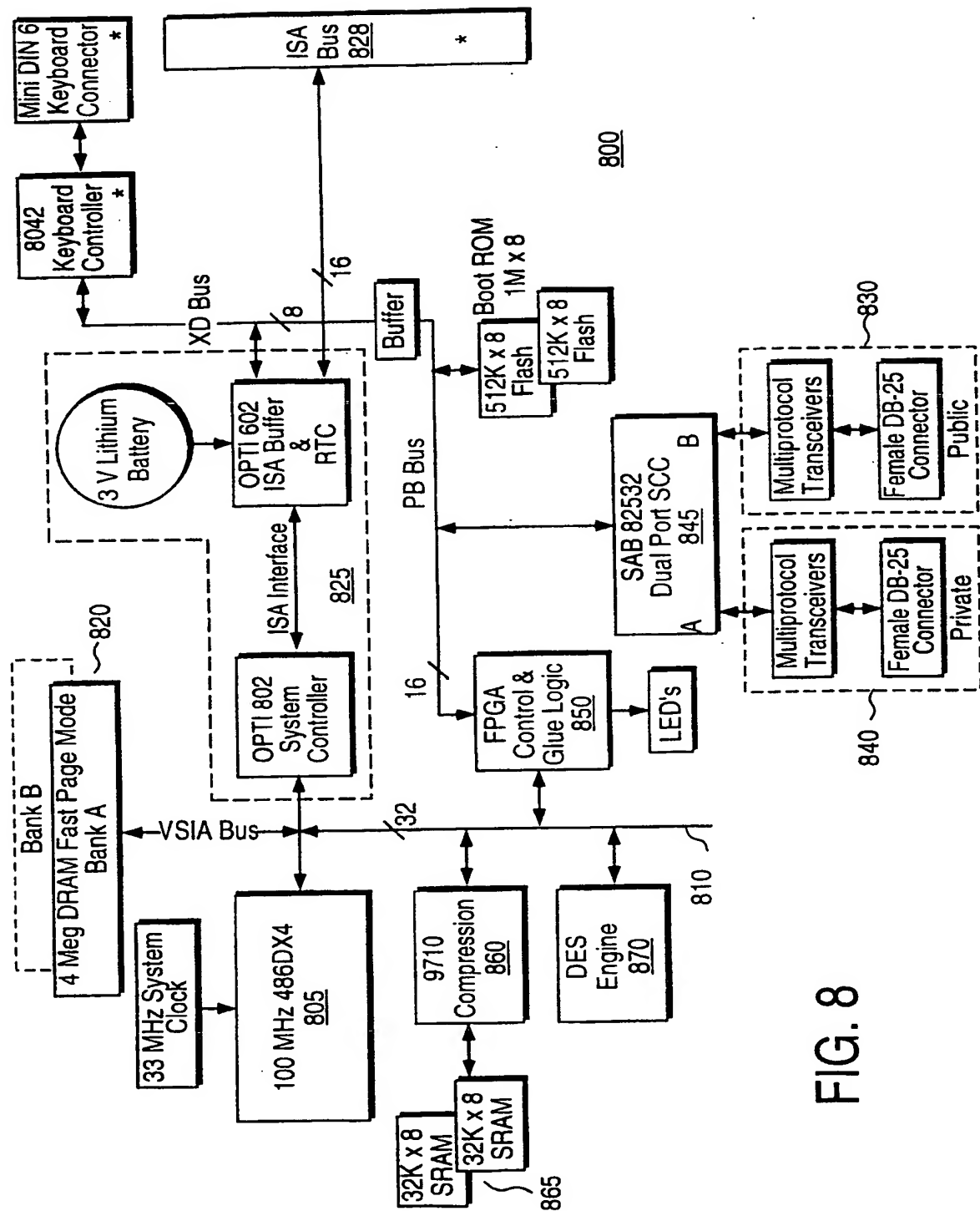


FIG. 8

INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 98/12226

A. CLASSIFICATION OF SUBJECT MATTER
IPC 6 H04L12/46 H04L29/06

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC 6 H04L

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

Electronic data base consulted during the international search (name of data base and, where practical, search terms used)

C. DOCUMENTS CONSIDERED TO BE RELEVANT

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| A | US 4 897 874 A (LIDINSKY WILLIAM P ET AL) 30 January 1990 see column 2, line 39 - line 55 see figure 15 --- -/- | 1-16 |

☒ Further documents are listed in the continuation of box C.

☒ Patent family members are listed in annex.

* Special categories of cited documents :

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Date of the actual completion of the international search

24 August 1998

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Name and mailing address of the ISA

European Patent Office, P.B. 5818 Patentlaan 2
NL - 2280 HV Rijswijk
Tel. (+31-70) 340-2040, Tx. 31 651 epo nl,
Fax: (+31-70) 340-3016

Authorized officer

Perez Perez, J

INTERNATIONAL SEARCH REPORT

International Application No

PCT/US 98/12226

C.(Continuation) DOCUMENTS CONSIDERED TO BE RELEVANT

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